



Latin American and Caribbean Internet Addresses Registry
Registro de Direcciones de Internet para **América Latina** y **Caribe**
Registro de Endereços da Internet para **América Latina** e **Caribe**

An Overview of DNSSEC

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DNSSEC

- Cryptography 101
- DNSSEC
- Where DNSSEC ?
- How does DNSSEC work ?
- New Resource Records
- Trust Chains



CRYPTOGRAPHY



Cryptography

- Cryptography concepts we'll need for DNSSEC
 - ◆ Public-key Cryptography
 - ◆ Hashing algorithms
 - ◆ Digital signatures
 - ◆ Trust Chains

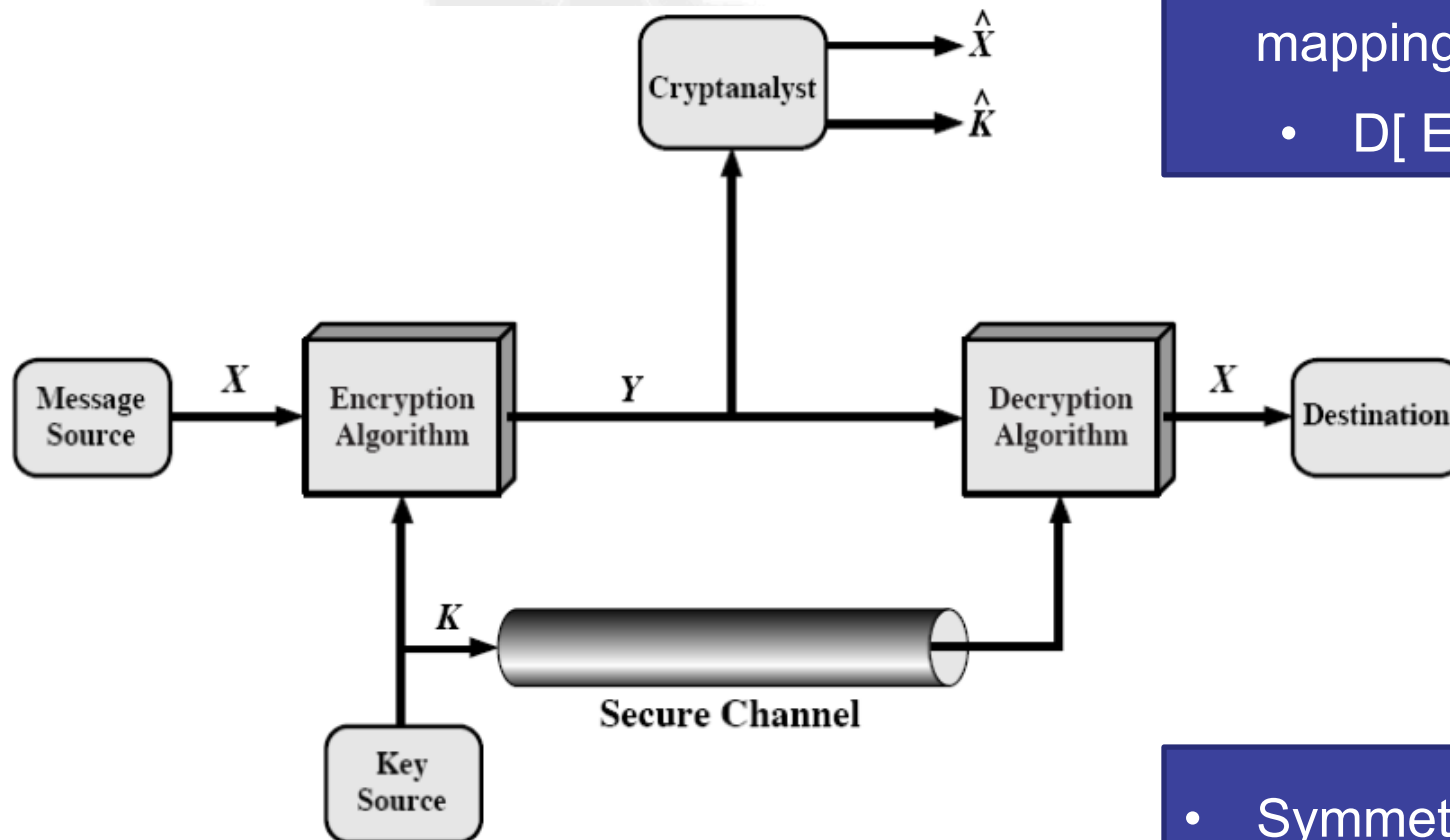


Cryptography (ii)

- Let's imagine two parties which need to communicate in a private manner. They will like to see certain **properties** enforced in their data exchanges.
 - ◆ They'd like to be sure that no one else has been **able to read** their messages (**privacy** property)
 - ◆ They'd like to be sure that no one else has been **able to change or alter** their messages (**integrity** property)
 - ◆ They'd like to be sure that the party who sends a message is really who it claims to be (**authentication** property)

Symmetric Cryptography

[Source: Stallings]



- $E[.]$ y $D[.]$ are two functions both which are inverse mappings of each other
- $D[E [X]] = X$

K (the key) is a *parameter* introduced to ease compromise recoveries

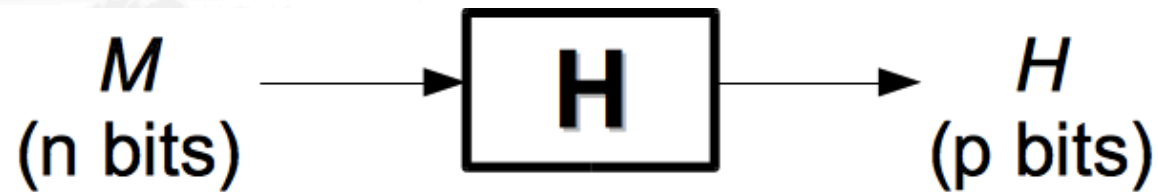
- Symmetric Cryptography
- $D_K[E_K [X]] = X$

Cryptographic Hashes



- H is a transformation with the following properties
 - ◆ $p \ll n$
 - ◆ For each algorithm "p" is a given value
 - ◆ $\text{len}(H)$ is fixed regardless of $\text{len}(M)$
- This means that **collisions** do exist
- Collision: If for a pair $M1$ and $M2$, $H(M1) == H(M2)$, then $M1$ and $M2$ represent a collision
- If $H()$ is chosen and designed carefully then finding collisions is very difficult

Cryptographic Hashes (ii)



- Intuitively
 - ◆ The more "random" the result of a hash "looks", the better it is
- Some well-known algorithms:
 - MD5
 - 128 bits
 - SHA1 / SHA256
 - 160 / 256 bits



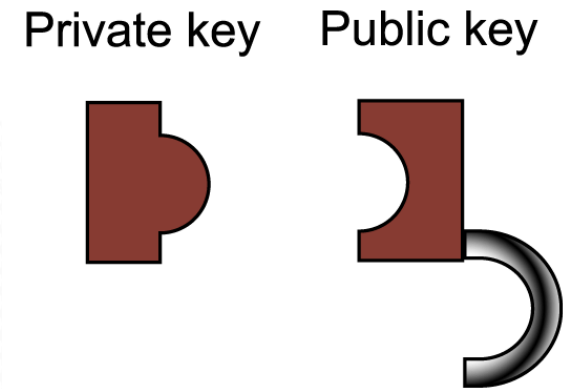
Public-key Cryptography

- Key distribution was always *the* weak point in traditional (symmetric) cryptography
- A lot effort was put to find workarounds and alternatives
- Breakthrough: (*Diffie-Hellman ca. 1976*) “Public-key Cryptography”
- A public-key cryptosystem has the following properties:
 - ◆ $D_{K_1}[E_{K_2} [X]] = X$
 - ◆ D cannot be easily found even if E is known
 - ◆ E cannot be broken with a chosen plaintext attack



Public-key Cryptography (ii)

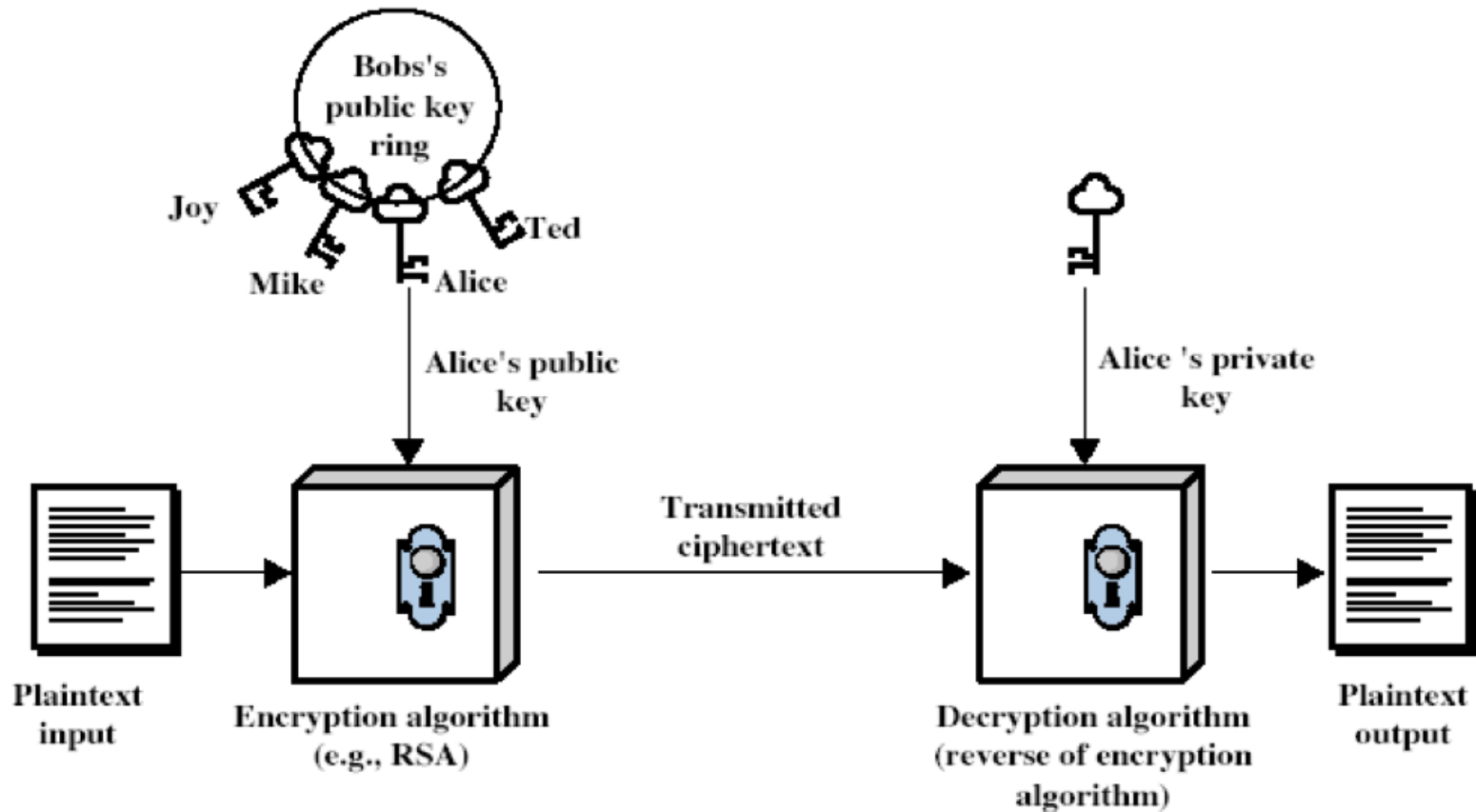
- Each party generates a keypair, that is one public and one secret key
 - ◆ K_{pub} , K_{priv}
 - ◆ Both keys in the pair are related
 - If one is given the other is also given
- When transmitting a message “X” from A -> B the following computation takes place:
 - ◆ $Y = E [K_{pub_B}, X]$
- When B receives the encrypted message the following computation takes place:



- ◆ $X' = D [K_{priv_B}, Y]$

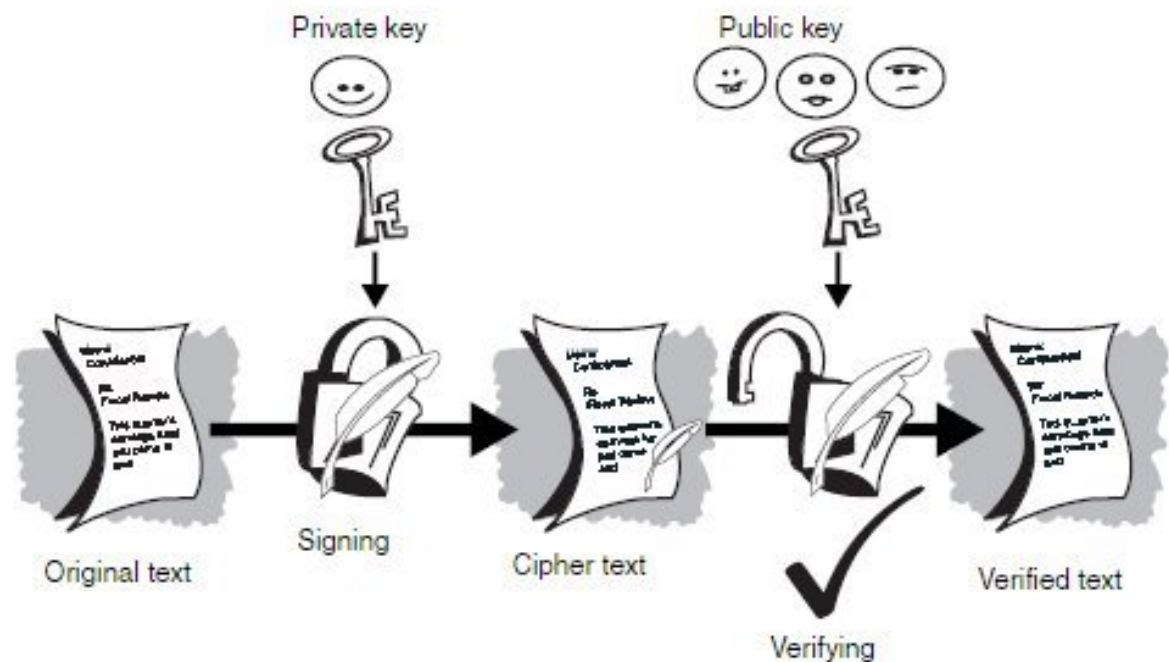
Public-key Cryptography (iii)

- [Source: Stallings]



Digital Signatures

- Goal:
 - ◆ *Create integrity proofs of digital documents*
- Usually implemented using public-key cryptography



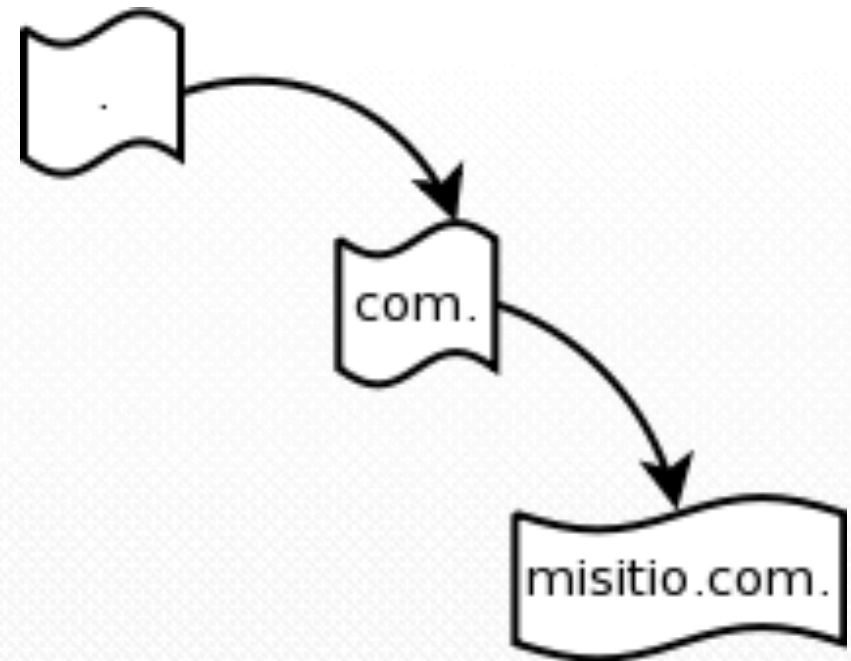


Digital Signatures (iii)

- Given M , a digital document to be signed by party A (lice) to be received by party B (ob)
 - ◆ A computes:
 - A hash for M , $H = \text{Hash}[M]$
 - A signature for M , $F = E[K_{\text{priv}_A}, H]$
 - ◆ A sends the pair $\{M, F\}$ to B
- When B receives the encrypted message the following computation takes place:
 - ◆ A hash for M , $H' = \text{Hash}[M]$
 - ◆ The hash of the signature is recalculated, $H = D[K_{\text{pub}_A}, F]$

Digital Signatures (iii)

- Trust Chains
 - ◆ Each level in the hierarchy signs data in the next one
 - ◆ The root needs to be analyzed separately
 - ◆ Validation can be either
 - Top down
 - Up down



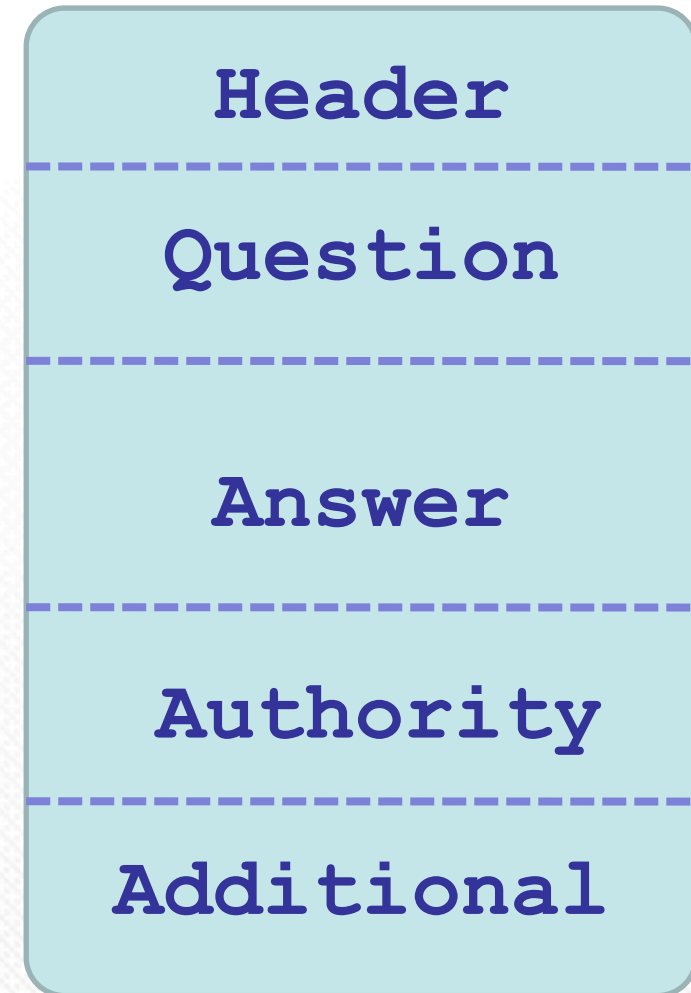


DNSSEC: MOTIVATION



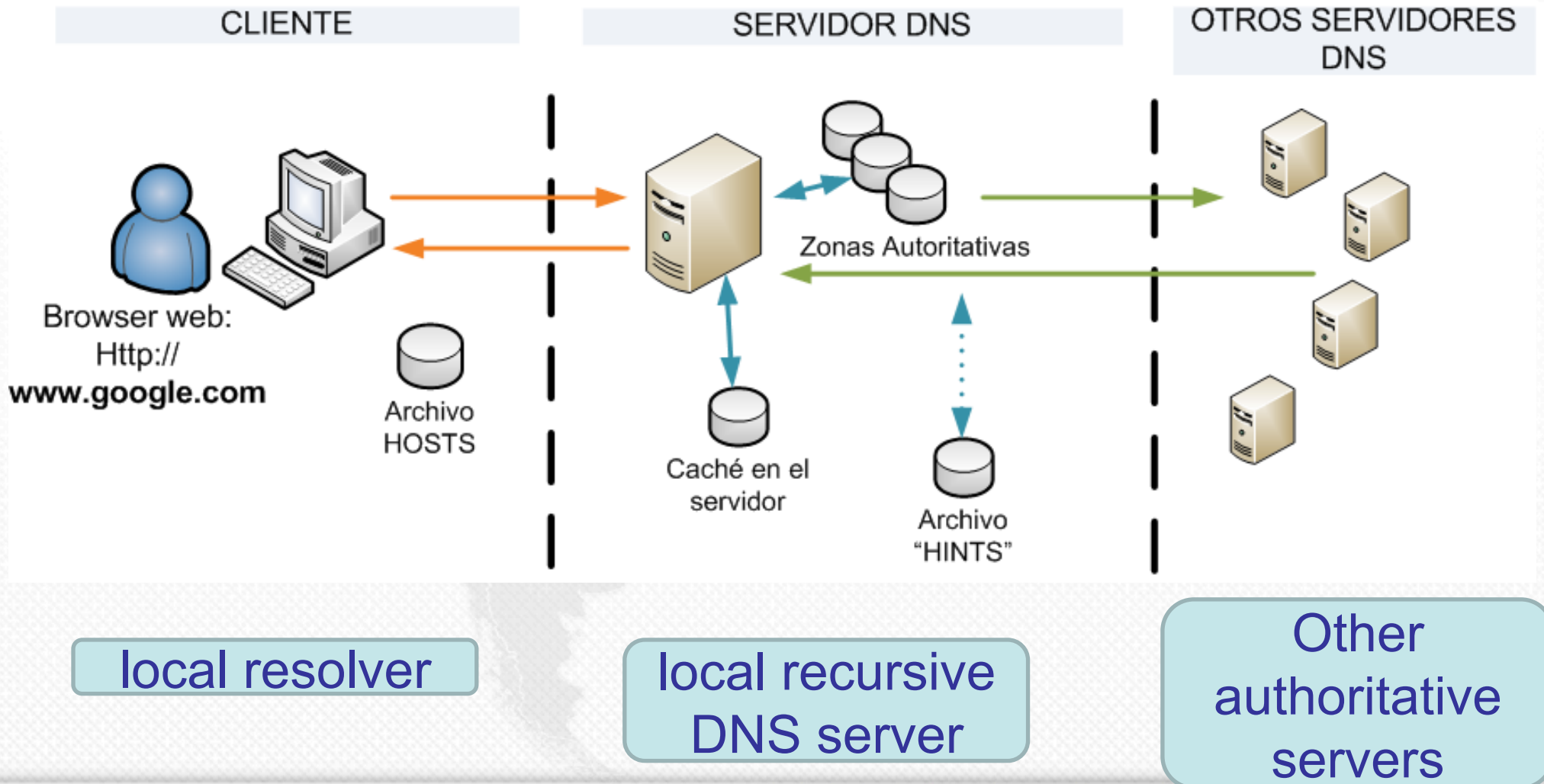
Protocol Specification

- Overview of DNS's wire packet format
 - ◆ Header
 - Protocol header
 - Flags (QR, RA, RD,...)
 - ◆ Question Section
 - Query we send to the DNS server
 - ◆ Tuples (*Name, Type, Class*)
 - ◆ Answer Section
 - RRs that answer the query (if any are available), also in (N, T, C) tuple format
 - ◆ Authority Section
 - RRs pointing to authoritative servers (optional)
 - ◆ Additional Section
 - RRs that may be useful to the querying client (according to the server answering the query)

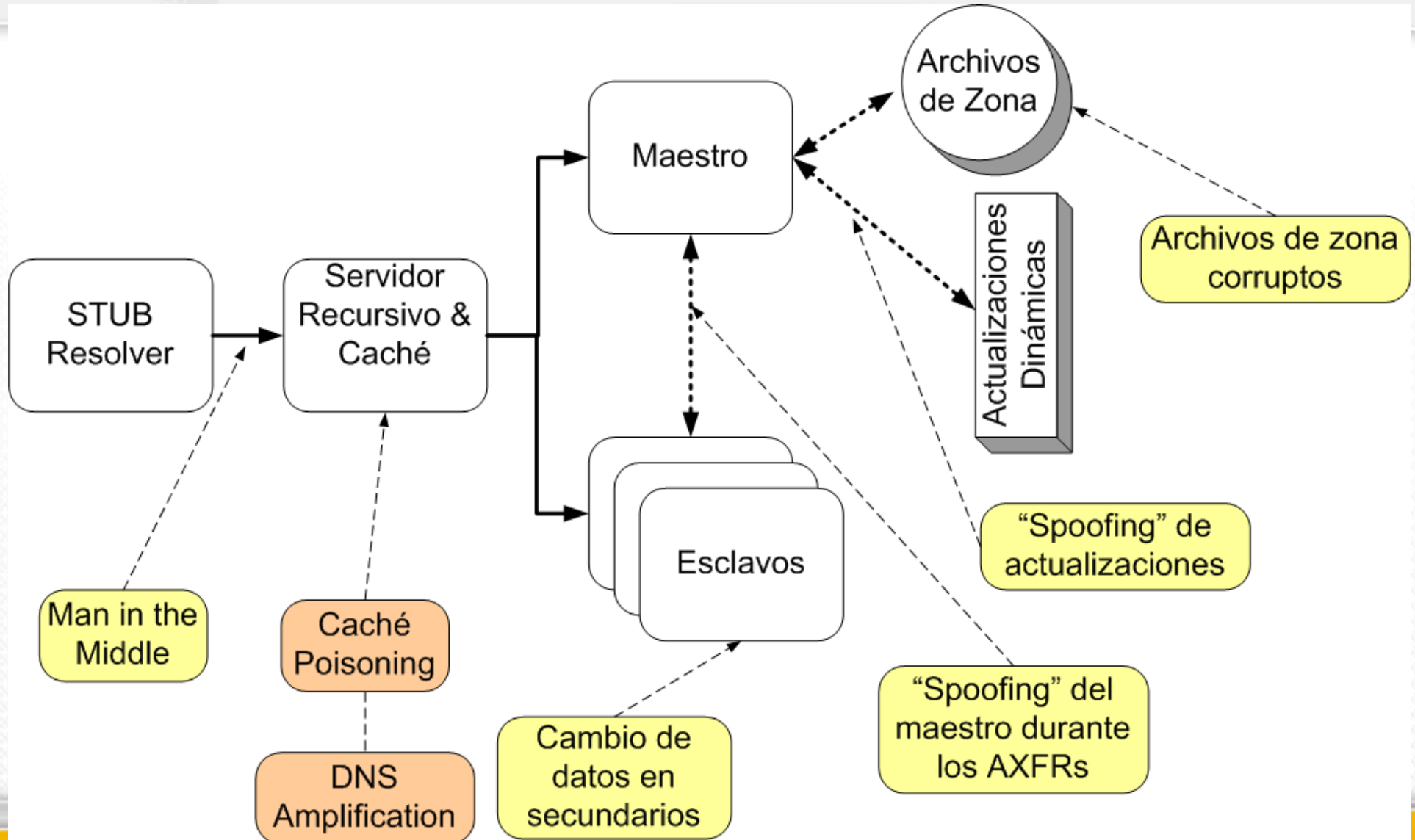




DNS Queries



Attack Vectors in DNS





Vulnerabilities in DNS

- DNS transmitted data is more prone to spoofing as it is mostly transported over UDP
 - ◆ Between master and slaves (AXFR)
 - ◆ Between masters and clients (AXFR) “*resolver*”
- Currently the DNS protocol does not have a way to validate information found in a query response
 - ◆ Vulnerable to different *poisoning* techniques
 - ◆ Poisoned data can be a problem for long periods of time depending on the TTL values of the zones
- Neither do slave servers have a way to authenticate the master servers they're talking to



Introducing DNSSEC

- Threat analysis in the DNS system
 - ◆ RFC 3833: “*Threat Analysis of the Domain Name System (DNS)*”
- DNSSEC:
 - ◆ “*DNS Security Extensions*”
 - ◆ RFC 4033, 4034, 4035
 - ◆ ~ May 2005



What does DNSSEC Protect us from?

- DNSSEC will protect us from data corruption and spoofing
 - ◆ It provides a way to validate both the integrity and the authenticity of the records contained in a DNS zone
 - **DNSKEY/RRSIG/NSEC**
 - ◆ It provides a way to delegate trust in public keys (trust chains)
 - **DS**
 - ◆ It provides a way to authenticate zone transfers between masters and slaves
 - **TSIG**



DNSSEC Introduction

- DNSSEC is **not** a new protocol
- Is a set of **extensions** to the DNS protocol as we know it
 - ◆ Changes to the wire protocol (EDNS0)
 - Maximum UDP query response extended from 512 to 4096 bytes
 - ◆ New resource records added
 - RRSIG, DNSKEY, DS, NSEC
 - ◆ New flags added
 - Checking Disabled (CD)
 - Authenticated Data (AD)



DNSSEC Introduction (2)

- New RRs
 - ◆ RRSIG: *Resource Record Signature*
 - ◆ DNSKEY: *DNS Public Key*
 - ◆ DS: *Delegation Signer*
 - ◆ NSEC: *Next Secure*
- New Flags:
 - ◆ AD: authenticated data
 - ◆ CD: checking disabled



DNSSEC Introduction (3)

- A resource record in DNS is a five-value tuple
 - ◆ (*name, class, type, TTL, value*)
- The record:
 - ◆ www.company.com. 86400 IN A 200.40.100.141
 - ◆ Is represented by the tuple:
 - Name (www.company.com)
 - Class (IN)
 - Type (A)
 - TTL (86400 seconds)
 - Value (200.40.100.141)



DNSSEC Introduction (4)

- ◆ *Resource Record Sets (RRSets)*
 - DNSSEC works by signing RRSets (not individual RRs)
 - An RRSet is a set of resource records that share the same:
 - ◆ Class
 - ◆ Type
 - ◆ Name
- ◆ Sample RRSet (TTL omitted for clarity)
 - `www IN A 200.40.241.100`
 - `www IN A 200.40.241.101`



DNSSEC Introduction (5) Zone Signing

- A key pair is created for each zone
 - ◆ Each zone has at least one key pair
 - ◆ The private key is kept, well, private
 - The private key is used to sign the RRsets in the zone
 - ◆ The public key is published in DNS using DNSKEY records
 - The private key is also used to verify the signatures of the RRsets
 - ◆ An RRSet can have multiple signatures generated using different key pairs



DNSSEC Introduction (6)

- The digital signature of a RRSet is returned in a special RRSIG record with the query answer
- Example:

```
~ carlosm$ dig +dnssec www.nic.se
```

```
;; flags: qr rd ra; QUERY: 1, ANSWER: 2, AUTHORITY: 4, ADDITIONAL: 1
```

```
;; ANSWER SECTION:
```

```
www.nic.se.          60      IN      A       212.247.7.218
www.nic.se.          60      IN      RRSIG   A 5 3 60 20101021132001
20101011132001 23369  nic.se. HeeUZ5h5iExK5uU1SuNRIf2Dbmh2/
aWV8FkjmzixUzTAVrHv39PfmfnG DHdHoZxoz85hqqYiWb
+t9EZh5+iaqxQk8AxRDic9Nn6Wxif0oWeS+IUKQ
rVyqXf1NtkZvu1A325vwa8obtbeVGVkhqg6bDIjKYeHixjlQ4cRoFcEW Izk=
```

```
:: AUTHORITY SECTION:
```

```
nic.se.             2974    IN      NS      ns3.nic.se.
nic.se.             2974    IN      NS      ns2.nic.se.
nic.se.             2974    IN      NS      ns.nic.se.
nic.se.             3600    IN      RRSIG   NS 5 2 3600
20101021132001 20101011132001 23369  nic.se. GSzAUC3SC3D0G/
iesCOPnVux8WkQx1dGbw491RatXz53b7SY0pQuyT1W
eb063Z62rtX7etynNcJwpKLYTG9FeMbDceD9af3KzTJHxq6B+Tpmmxyk
FoKAVaV0cHTcGUXS0bFquGr5/03G79C/YHJmXw0bHun5ER5yr0t0LegU IAU=
```




Trust Chains

- How do clients verify a zone's RRsets?
 - ◆ It queries for the corresponding DNSKEY
 - ◆ The necessary computations are carried out and then compared with the signature in the RRSIG
 - If they match the signatures are valid
- But, how can we trust the DNSKEY? It listed on the same zone we want to verify!
 - ◆ We need to validate the **trust chain**



Trust Chains (ii)

- DS Record “*Delegation Signature*”
 - ◆ DS records "sign" the keys in their child zones
 - ◆ In this way one can also verify the DNSKEY as it is signed when the parent zone is signed
- DS records contain a hash of the public key
 - ◆ That is a hash of the DNSKEY's record content
- DS records in the parent zone are signed with the keys of the parent zone
- To complete the full trust chain we also need the **root of the DNS** to be signed



Trust Chains (iii)

- What about the root zone ?
 - ◆ The root zone has no parent zone where a DS record could be placed
 - ◆ The DNS root has been signed since July 2010
 - [<http://www.root-dnssec.org>]
 - ◆ The DS record for "." is obtained out-of-band and installed locally in each server
 - [<http://data.iana.org/root-anchors/root-anchors.xml>]
 - . IN DS
49AAC11D7B6F6446702E54A1607371607A1A41855200F
D2CE1CDDE32F24E8FB5



DNSSEC Introduction (9) Root Zone Signing

- How is the the root trust anchor verified?
- It is verified also out-of-band
 - ◆ It can be downloaded using http/https
 - ◆ Several validation mechanisms are in place (X.509 certs, PGP signatures)
 - ◆ It is locally installed in the same way the root zone itself is configured locally



DNSSEC Introduction (10) Denial of Existence

- "NXDOMAIN" answers
 - ◆ Provide “denial of existence” answers via a flag on the “Header” pseudo-section
 - ◆ NXDOMAINS are cached in the same way as other responses are
 - ◆ Forging NXDOMAINs is a DDoS attack vector
- How can non-existence be signed ?
 - ◆ We need an RRSet to sign
 - Remember that DNSSEC always signs RRsets
 - ◆ Two different techniques have been proposed:
 - NSEC and NSEC3



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Thank You!

Questions?